

Eliminations and echelon forms in exact linear algebra

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University of Waterloo, ON, Canada,
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Gaussian elimination in Computer Algebra

Linear system solving: over $\mathbb{Z}_p, \mathbb{Z}, \mathbb{Q}$ (Crypto, Number Theory)

Polynomial system solving: Gröbner basis (Robotics, Crypto)

Linear dependencies: rank, basis (of vector spaces, free modules,
Krylov spaces, ...) (Number Theory)

Determinants: certificate of similarity (Graph Theory)

Designing efficient dense gaussian elimination routines over an exact ring/field.

- Extensively studied for numerical computations
- Specificities of exact computations:
 - ▶ No partial/full pivoting
 - ▶ Rank profile matters
- size of coefficients (e.g. compressed in $\text{GF}(2)$) \Rightarrow asymmetry

Study originating from, the design of the libraries

- FFLAS-FFPACK: word size finite fields
- M4RI: $\text{GF}(2)$

1 Decompositions and factorizations

- Gaussian elimination based matrix decompositions
- Relations between decompositions

2 Algorithms

- Block recursive gaussian elimination
- Memory allocations
- Time complexity

3 Parallelization

4 Algorithms into practice: the case of GF(2)

Outline

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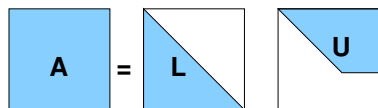
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LU decomposition

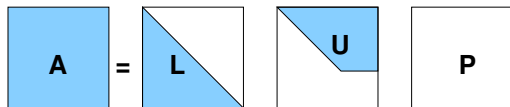


- L unit lower triangular,
- U non-sing upper triangular

Exists for

- matrices having the generic rank profile (every leading principal minor is non zero)

LUP, PLU decomposition

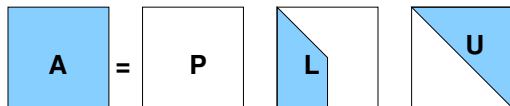
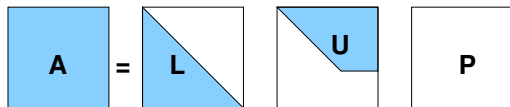


- P a permutation matrix

Exists for

- Any non-singular matrix
- Or any matrix with generic row rank profile

LUP, PLU decomposition

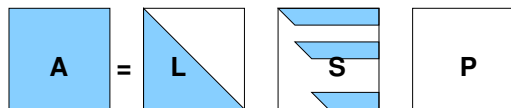


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LSP, LQUP, PLUQ decompositions

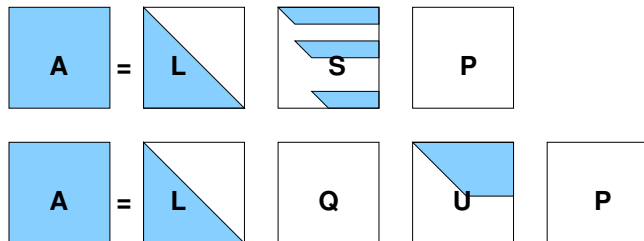


- S : semi-upper triangular,
- Q permutation matrix

Exists for

- any $m \times n$ matrix

LSP, LQUP, PLUQ decompositions

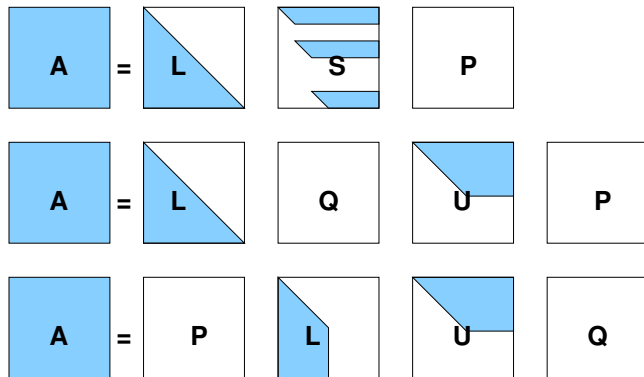


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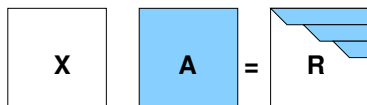
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Echelon form decomposition

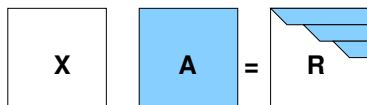
Row Echelon Form $XA = R$



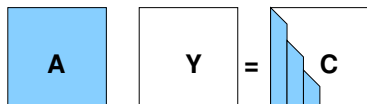
- X, Y : non-singular transformation matrices
- R, C : matrices in row/col echelon form

Echelon form decomposition

Row Echelon Form $XA = R$



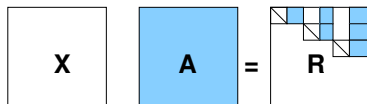
Column Echelon Form $AY = C$



- X, Y : non-singular transformation matrices
- R, C : matrices in row/col echelon form

Reduced echelon form decomposition

Row Reduced Echelon Form $XA = R$


$$X A = R$$

- X, Y : non-singular transformation matrices
- R, C : matrices in reduced row/col echelon form

Reduced echelon form decomposition

Row Reduced Echelon Form $XA = R$

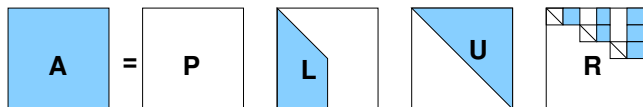
The diagram shows three square matrices. The first matrix, labeled **X**, is white. The second matrix, labeled **A**, is light blue. An equals sign follows. The third matrix, labeled **R**, is white with a blue upper triangular pattern, where the diagonal and the upper right portion are shaded blue.

Column Reduced Echelon Form $AY = C$

The diagram shows three square matrices. The first matrix, labeled **A**, is light blue. The second matrix, labeled **Y**, is white. An equals sign follows. The third matrix, labeled **C**, is white with a blue lower triangular pattern, where the diagonal and the lower left portion are shaded blue.

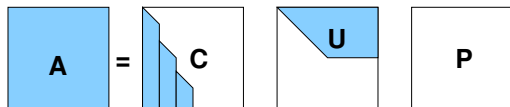
- X, Y : non-singular transformation matrices
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Turing factorization



- PLU decomposition of X^{-1} , the inverse of the transformation matrix to REF

CUP and PLE decompositions

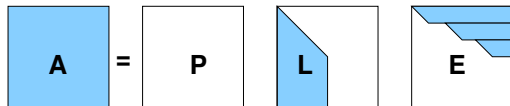
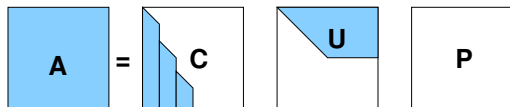


- C : column echelon form
- E : row echelon form

Exists for

- any $m \times n$ matrix

CUP and PLE decompositions



- C: column echelon form
- E: row echelon form

Exists for

- any $m \times n$ matrix

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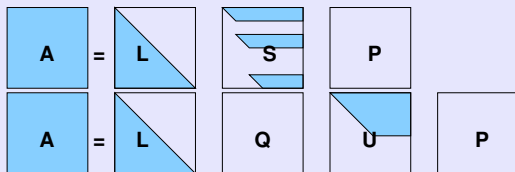
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Relations: up to permutations

From LSP to LQUP

$$S = QU$$



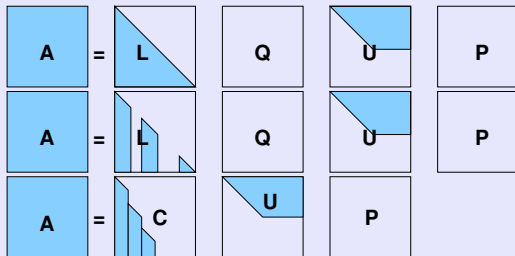
Fact

The first $r = \text{rank}(A)$ values of the permutation Q are monotonically increasing.

Relations: up to permutations

From LQUP to CUP

$$C = LQ$$

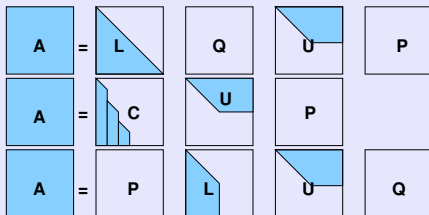


And to PLE using transposition: $PLE(A^T) = CUP(A)^T$

Relations:

From LQUP to PLUQ

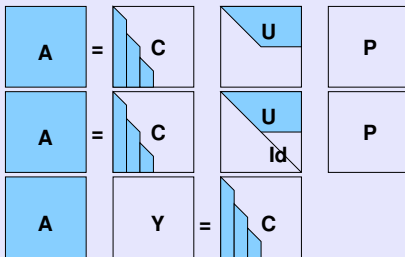
$$P \leftrightarrow Q, L \leftarrow Q^T L Q$$



Relations:

From CUP to ColumnEchelon form

$$Y = P^T \begin{bmatrix} U \\ I_{n-r} \end{bmatrix}^{-1}$$

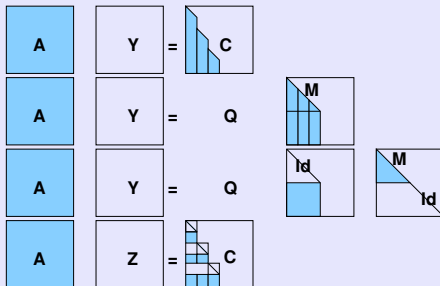


Similarly, from PLE to RowEchelon form

Relations:

From Column Echelon form to Reduced Column Echelon form

$$Z = Y \begin{bmatrix} M \\ I_{n-r} \end{bmatrix}^{-1}$$

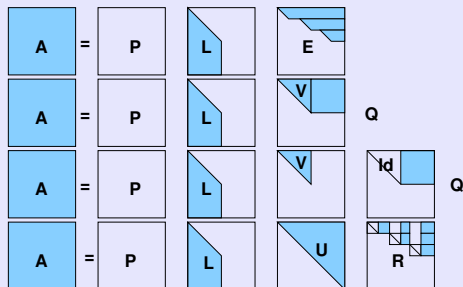


Similarly, from PLE to RowEchelon form

Relations:

From PLE to Turing

$$U = V,$$
$$R = \begin{bmatrix} V & \\ & I_{n-r} \end{bmatrix}^{-1} E$$



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Algorithms: main types

Three ways to group operations:

- ① simple iterative
 - ▶ Apply the standard Gaussian elimination in dimension n
 - ▶ Main loop **for** $i=1$ to n
- ② block algorithms
 - ① block iterative (Tile)
 - ★ Apply Gaussian elimination in dimension n/k over blocks of size k
 - ★ Main loop: **for** $i=1$ to n/k
 - ② block recursive
 - ★ Apply Gaussian elimination in dimension 2 recursively on blocks of size $n/2^i$
 - ★ Main loop: **for** $i=1$ to 2

Type of algorithms

Data locality: prefer block algorithms

- cache aware: block iterative
- cache oblivious: block recursive

Base case efficiency: simple iterative

Asymptotic time complexity: block recursive

Parallelization: block iterative

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Block recursive gaussian elimination

Author	Year	Computation	Requirement
Strassen	69	Inverse	gen. rank prof.
Bunch, Hopcroft	74	LUP	gen. row rank prof.
Ibarra, Moran, Hui	82	LSP, LQUP	none
Schönage, Keller-Gerig	85	StepForm	none
Storjohann	00	Echelon, RedEch	none
here	11	CUP,PLE,PLUQ	none

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Comparison according to

- No requirement on the input matrix

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Comparison according to

- No requirement on the input matrix
- Rank sensitive complexity

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Comparison according to

- No requirement on the input matrix
- Rank sensitive complexity
- Memory allocations
- Constant factor in the time complexity

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- **Memory allocations**
- Time complexity

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Memory requirements:

Definition

In place = *output overrides the input and computation does not need extra memory (considering Matrix multiplication $C \leftarrow C + AB$ as a black box)*

Remark: a unit lower triangular and an upper triangular matrix can be stored on the same $m \times n$ storage!

TRSM: TRIangular Solve with Matrix

$$\begin{bmatrix} A & B \\ C & \end{bmatrix}^{-1} \begin{bmatrix} D \\ E \end{bmatrix} = \begin{bmatrix} A^{-1} & \\ & I \end{bmatrix} \begin{bmatrix} I & -B \\ & I \end{bmatrix} \begin{bmatrix} I & \\ & C^{-1} \end{bmatrix} \begin{bmatrix} D \\ E \end{bmatrix}$$

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Compute $F = C^{-1}E$

(Recursive call)

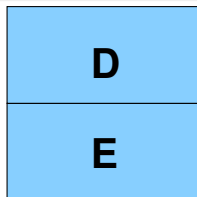
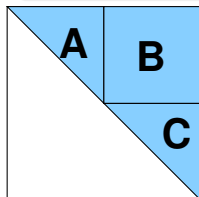
Compute $G = D - BF$

(MM)

Compute $H = A^{-1}G$

(Recursive call)

Return $\begin{bmatrix} H \\ F \end{bmatrix}$



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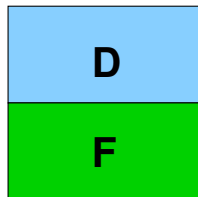
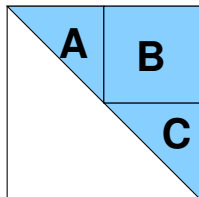
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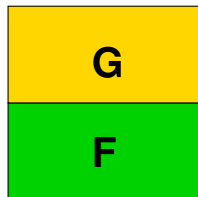
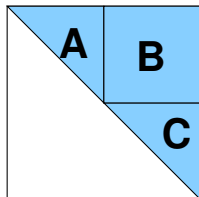
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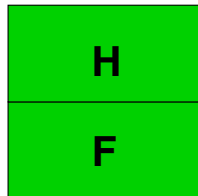
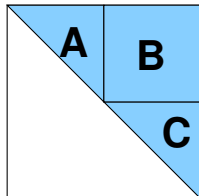
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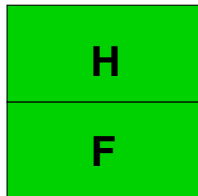
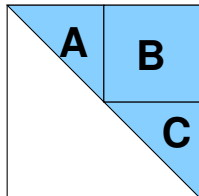
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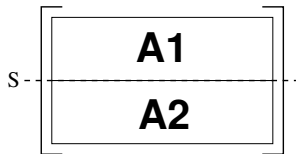
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Return $\begin{bmatrix} H \\ F \end{bmatrix}$

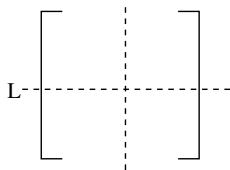


- $\mathcal{O}(n^{\omega})$
- In place

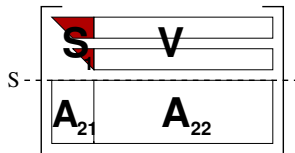
The LSP algorithm



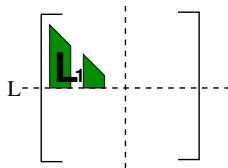
1 Split A Row-wise



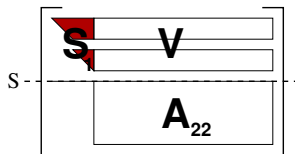
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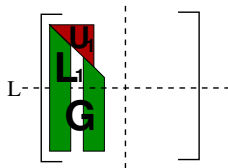
- 1 Split A Row-wise
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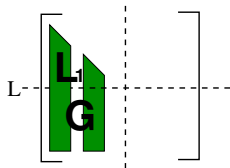
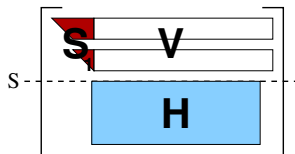
The LSP algorithm



- 1 Split A Row-wise
- 2 Recursive call on A_1
- 3 $G \leftarrow A_{21} U_1^{-1}$ (trsm)

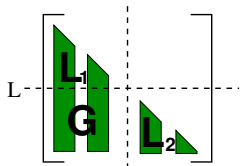
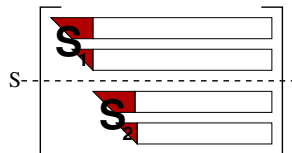


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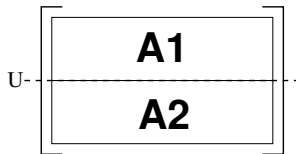
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- 4 $H \leftarrow A_{22} - G \times V$ (MM)

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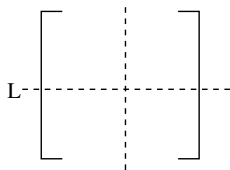


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- 3 $G \leftarrow A_{21} U_1^{-1}$ (trsm)
- 4 $H \leftarrow A_{22} - G \times V$ (MM)
- 5 Recursive call on H

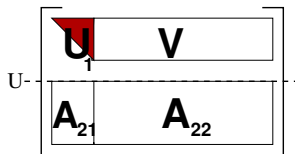
The LQUP algorithm



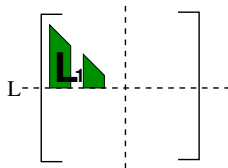
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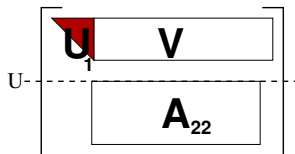
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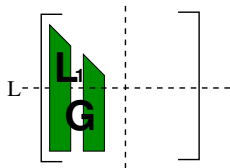
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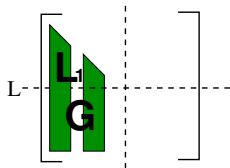
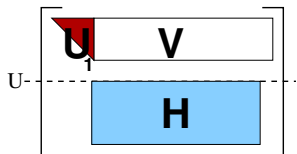
The LQUP algorithm



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- 3 $G \leftarrow A_{21} U_1^{-1}$ (trsm)

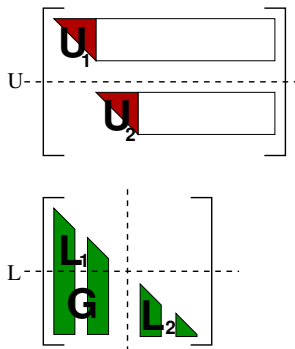


The LQUP algorithm



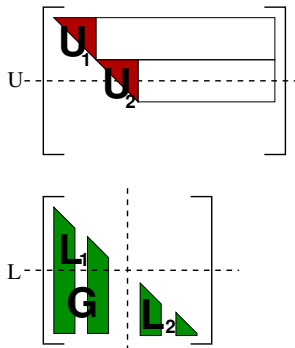
- 1 Split A Row-wise
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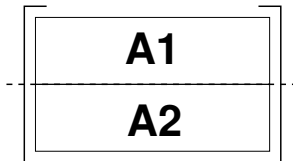
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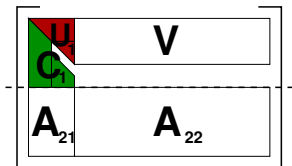
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The CUP decomposition

- 1 Split A Row-wise

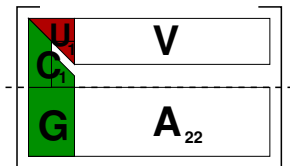


The CUP decomposition



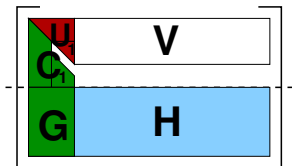
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The CUP decomposition



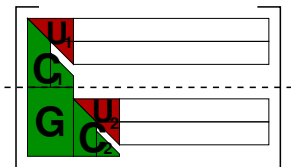
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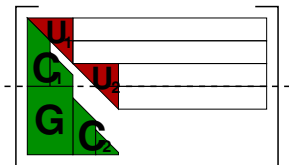
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The CUP decomposition



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- 4 $H \leftarrow A_{22} - G \times V$ (MM)
- 5 Recursive call on H
- 6 Row permutations

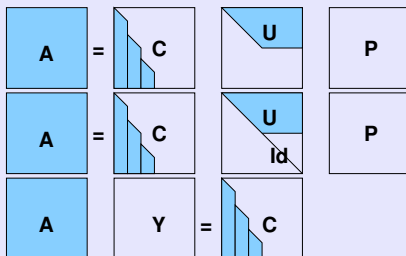
Memory: LSP vs LQUP vs PLUQ vs CUP

Decomposition	In place
LSP	N
LQUP	N
PLUQ	Y
CUP	Y

Echelon forms

From CUP to ColumnEchelon form

$$\begin{aligned} Y &= P^T \begin{bmatrix} U_1 & U_2 \\ & I_{n-r} \end{bmatrix}^{-1} \\ &= P^T \begin{bmatrix} U_1^{-1} & -U_1^{-1}U_2 \\ & I_{n-r} \end{bmatrix} \end{aligned}$$



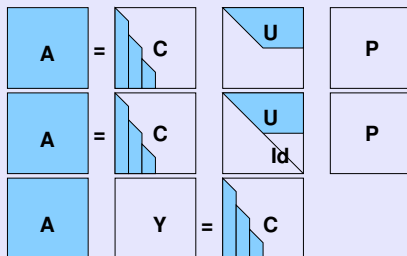
Additional operations:

$-U^{-1}U_2$ `trsm` (triangular system solve) **in-place**

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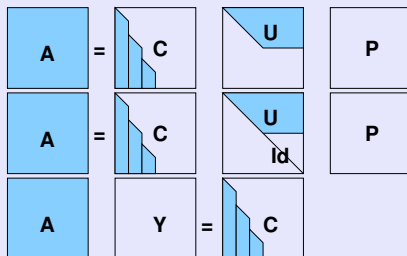
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From LQUP to Column Echelon

TRTRI: triangular inverse

$$\begin{bmatrix} U_1 & U_2 \\ & U_3 \end{bmatrix}^{-1} = \begin{bmatrix} U_1^{-1} & -U_1^{-1}U_2U_3^{-1} \\ & U_3^{-1} \end{bmatrix}$$

1: **if** $n = 1$ **then**

2: $U \leftarrow U^{-1}$

3: **else**

4: $U_2 \leftarrow U_3^{-1}U_2$

TRSM

5: $U_2 \leftarrow -U_2U_3^{-1}$

TRSM

6: $U_1 \leftarrow U_1^{-1}$

TRTRI

7: $U_3 \leftarrow U_3^{-1}$

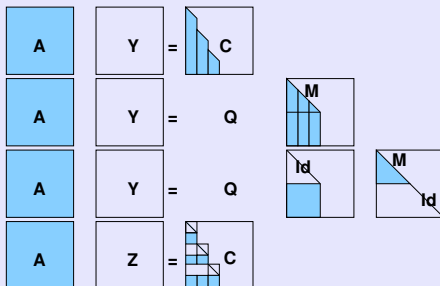
TRTRI

8: **end if**

Reduced Echelon forms

From Column Echelon form to Reduced Column Echelon form

$$Z = Y \begin{bmatrix} M & \\ & I_{n-r} \end{bmatrix}^{-1}$$



Similarly, from PLE to RowEchelon form

Again reduces to:

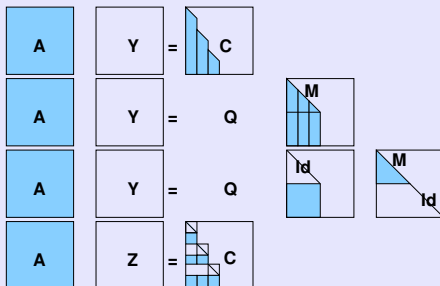
$U^{-1}X$: TRSM, **in-place**

U^{-1} : TRTRI, **in-place**

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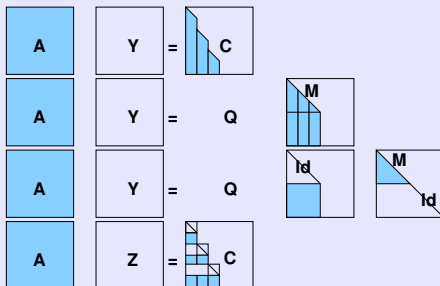
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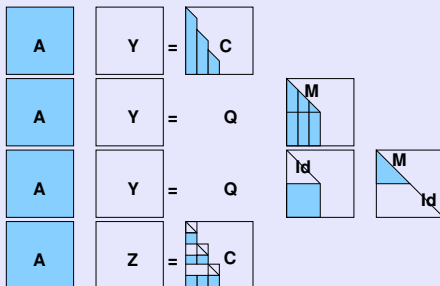
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From Echelon to Reduced Echelon

TRTRM: triangular triangular multiplication

$$\begin{bmatrix} U_1 & U_2 \\ & U_3 \end{bmatrix} \begin{bmatrix} L_1 & \\ L_2 & L_3 \end{bmatrix} = \begin{bmatrix} U_1 L_1 + U_2 L_2 & U_2 L_3 \\ & U_3 L_3 \end{bmatrix}$$

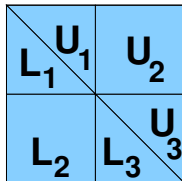
- | | |
|-----------------------------------|-------|
| 1: $X_1 \leftarrow U_1 L_1$ | TRTRM |
| 2: $X_1 \leftarrow X_1 + U_2 L_2$ | MM |
| 3: $X_2 \leftarrow U_2 L_3$ | TRMM |
| 4: $X_3 \leftarrow U_3 L_2$ | TRMM |
| 5: $X_4 \leftarrow U_3 L_3$ | TRTRM |

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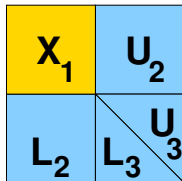


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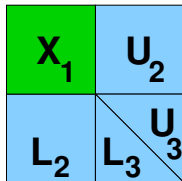


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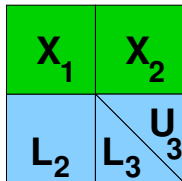


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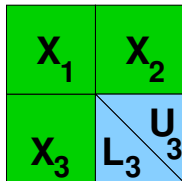


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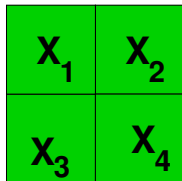


From Echelon to Reduced Echelon

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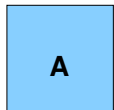
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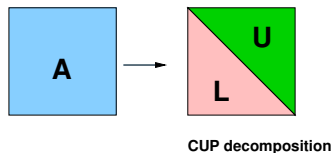


- $\mathcal{O}(n^\omega)$
- In place

Example: in place matrix inversion

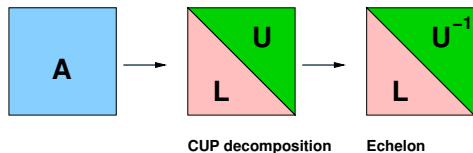


Example: in place matrix inversion



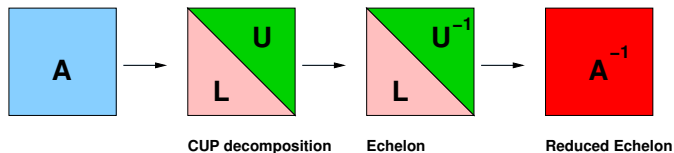
$$A = LU$$

Example: in place matrix inversion



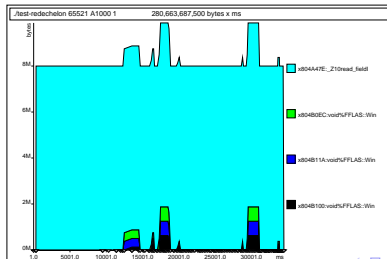
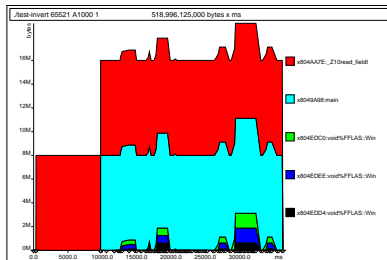
$$AU^{-1} = L$$

Example: in place matrix inversion



$$A(U^{-1}L^{-1}) = I$$

Experiments



Direct computation of the Reduced Echelon form

- Strassen 69: inverse of generic matrices
- Storjohann 00: Gauss-Jordan generalization for any rank profile

Matrix Inversion [Strassen 69]

$$\begin{bmatrix} A & B \\ C & D \end{bmatrix}^{-1} = \begin{bmatrix} A^{-1} & \\ & I \end{bmatrix} \begin{bmatrix} I & -B \\ & I \end{bmatrix} \begin{bmatrix} I & \\ & (D - CA^{-1}B)^{-1} \end{bmatrix} \begin{bmatrix} I & \\ CA^{-1} & I \end{bmatrix}$$

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- 1: Compute $E = A^{-1}$ (Recursive call)
- 2: Compute $F = D - CEB$ (MM)
- 3: Compute $G = F^{-1}$ (Recursive call)
- 4: Compute $H = -EB$ (MM)
- 5: Compute $J = HG$ (MM)
- 6: Compute $K = CE$ (MM)
- 7: Compute $L = E + JK$ (MM)
- 8: Compute $M = GK$ (MM)
- 9: Return $\begin{bmatrix} E & J \\ M & G \end{bmatrix}$

Strassen-Storjohann's Gauss-Jordan elimination

Problem

Needs to perform operations of the form $A \leftarrow AB$

\Rightarrow not doable in place by a usual matrix multiplication algorithm

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Workaround [Storjohann]:

- | | | |
|---|--------------------|------|
| 1 | Decompose $B = LU$ | LU |
| 2 | $A \leftarrow AL$ | trmm |
| 3 | $A \leftarrow AU$ | trmm |

Outline

1 Decompositions and factorizations

- Gaussian elimination based matrix decompositions
- Relations between decompositions

2 Algorithms

- Block recursive gaussian elimination
- Memory allocations
- **Time complexity**

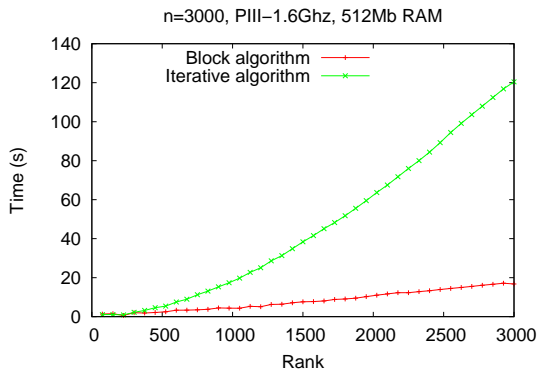
3 Parallelization

4 Algorithms into practice: the case of GF(2)

Rank sensitive time complexity

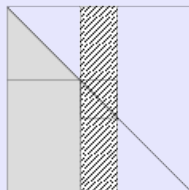
Fact

Algorithms *LSP*, *CUP*, *LQUP*, *PLUQ*, ... have a rank sensitive computation time: $\mathcal{O}(mnr^{\omega-2})$

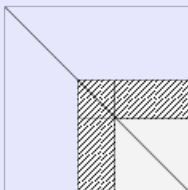


Rank sensitivity and Left/Right/Crout Looking variants

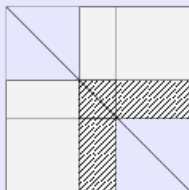
Right looking / Left looking/ Crout variants



Left-looking LU



Right-looking LU



Crout LU

- does not affect 2x2 splitting
- for block iterative: always the same rank sensitive complexity:

$$2n^2r - 2nr^2 + 2/3r^3$$

Time complexity: comparing constants

$$\mathcal{O}(n^\omega) = C_\omega n^3$$

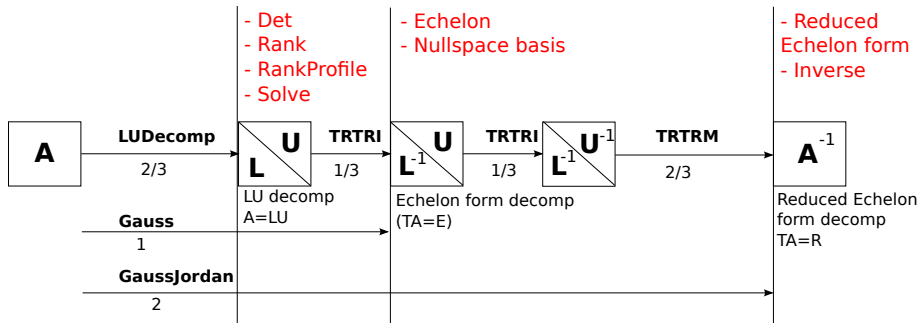
Algorithm	Constant C_ω	C_3	$C_{\log_2 7}$	in-place
MM	C_ω	2	6	×
TRSM	$\frac{C_\omega}{2^{\omega-1}-2}$	1	4	✓
TRTRI	$\frac{C_\omega}{(2^{\omega-1}-2)(2^{\omega-1}-1)}$	$\frac{1}{3} \approx 0.33$	$\frac{8}{5} = 1.6$	✓
TRTRM, CUP PLUQ LQUP,	$\frac{C_\omega}{2^{\omega-1}-2} - \frac{C_\omega}{2^{\omega-2}}$	$\frac{2}{3} \approx 0.66$	$\frac{14}{5} = 2.8$	✓
Echelon	$\frac{C_\omega}{2^{\omega-2}-1} - \frac{3C_\omega}{2^{\omega-2}}$	1	$\frac{22}{5} \approx 4.4$	✓
RedEchelon	$\frac{C_\omega(2^{\omega-1}+2)}{(2^{\omega-1}-2)(2^{\omega-1}-1)}$	2	$\frac{44}{5} = 8.8$	✓
StepForm	$\frac{5C_\omega}{2^{\omega-1}-1} + \frac{C_\omega}{(2^{\omega-1}-1)(2^{\omega-2}-1)}$	4	$\frac{76}{5} = 15.2$	×
GJ*	$\frac{C_\omega}{2^{\omega-2}-1}$	2	8	×

*: GJ: GaussJordan alg of [Storjohann00] computing the reduced echelon form

Applications to standard linalg problems

Problem	Using	C_ω	C_3	$C_{\log_2 7}$	In place
Rank					
RankProfile	GJ	$\frac{C_\omega}{2^{\omega-2}-1}$	2	8	×
IsSingular	CUP	$\frac{C_\omega}{2^{\omega-1}-2} - \frac{C_\omega}{2^{\omega-2}}$	0.66	2.8	✓
Det					
Solve					
Inverse	GJ	$\frac{C_\omega}{2^{\omega-2}-1}$	2	8	×
	CUP	$\frac{C_\omega(2^{\omega-1}+2)}{(2^{\omega-1}-2)(2^{\omega-1}-1)}$	2	8.8	✓

Summary



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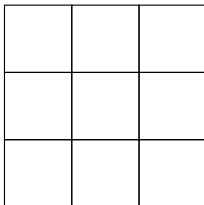
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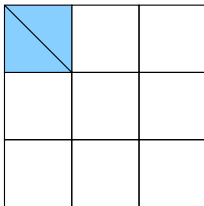
Parallelization

- Using block iterative algorithm
- Parallelizing the matrix multiplication updates
- Always a critical path of n on the diagonal



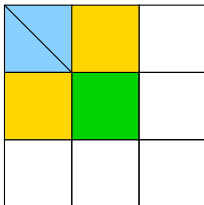
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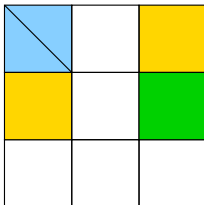
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- Parallelizing the matrix multiplication updates
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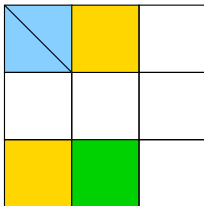
Parallelization

- Using block iterative algorithm
- Parallelizing the matrix multiplication updates
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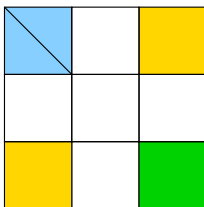
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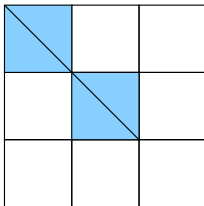
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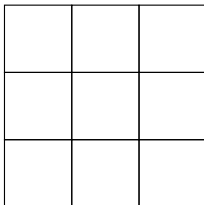
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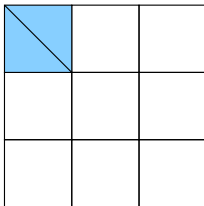
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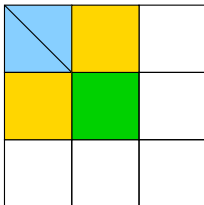
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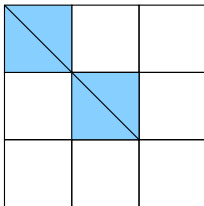
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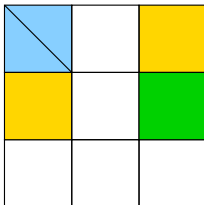
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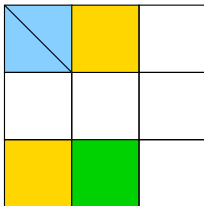
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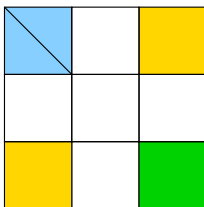
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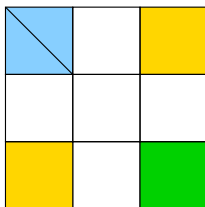
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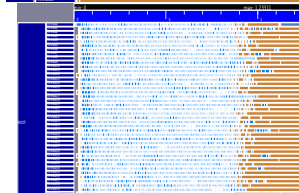
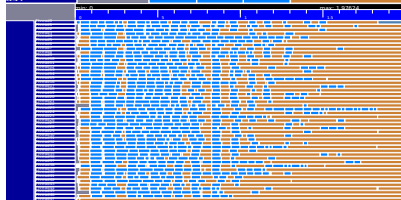
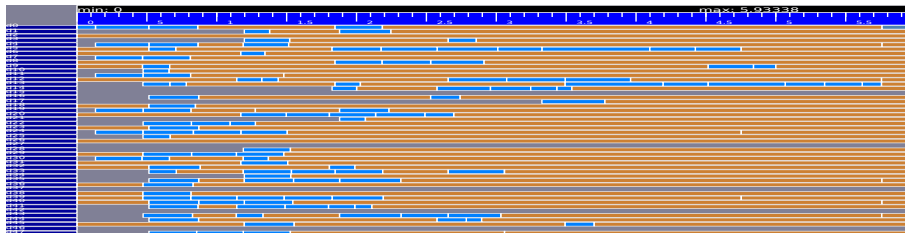
Scheduling tasks

- DAG, static scheduling, ...
- Dynamic scheduling, work-stealing

Advantage of exact computations:

- No need to update pivots (local pivoting) \Rightarrow less ops
- Multi-frontal approach made easier

Dynamic scheduling with Kaapi



Outline

1 Decompositions and factorizations

- Gaussian elimination based matrix decompositions
- Relations between decompositions

2 Algorithms

- Block recursive gaussian elimination
- Memory allocations
- Time complexity

3 Parallelization

4 Algorithms into practice: the case of GF(2)

The M4RI library [M Albrecht & Al.]:

Dense basic linear algebra over GF(2)

- Packed representation: `unsigned long long` is a vector of 64 coefficients
- Matrix multiplication:
 - ▶ Gray code table lookup (Methode of the 4 russians)
 $\Rightarrow \mathcal{O}(n^3 / \log n)$
 - ▶ Strassen on the coarse grain.
 $\Rightarrow \mathcal{O}((n/k)^\omega k^3 / \log k) = \mathcal{O}(n^\omega)$

Design of the gaussian elimination routine

PLE vs CUP

- Row major storage \Rightarrow easier to permute pivots along columns
 \Rightarrow PLE rather than CUP

PLE vs PLUQ

- PLUQ involves more back and forth pivoting
- Compact LAPACK representation of permutation: product of transpositions
 - ▶ Not possible to maintain for P in $\mathcal{O}(1)$ PLUQ \Rightarrow PLE rather than PLUQ

Design of the gaussian elimination routine

Structure of the algorithm

- block recursive PLE on the coarse grain ($\mathcal{O}(n^\omega)$)
- block iterative PLE with block size $k = \log n$
- simple iterative PLE on cache fitting blocks

Additional tricks:

- With gray code tables: inverting means reverse table look-up
⇒ TRSM as efficient as MM
- ...

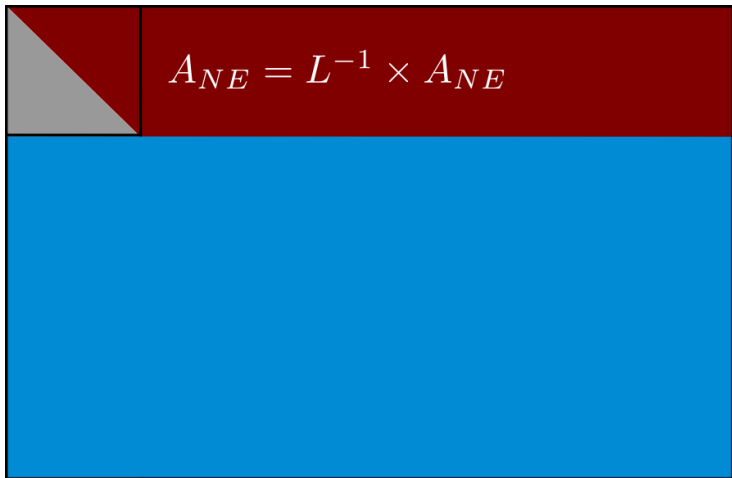


Visualisation

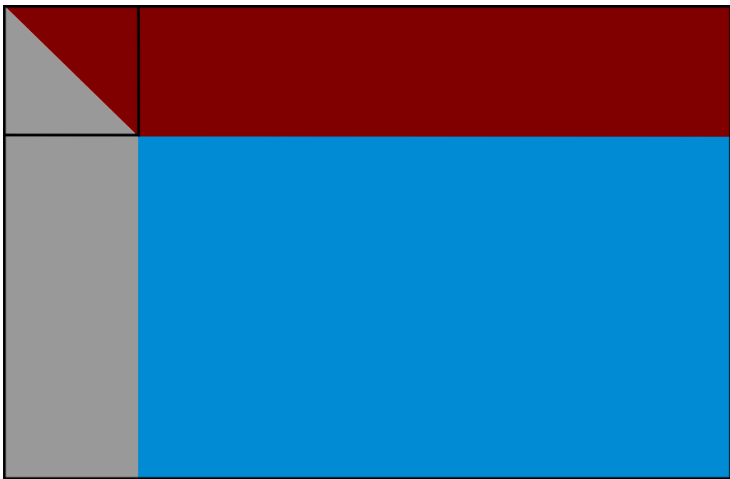


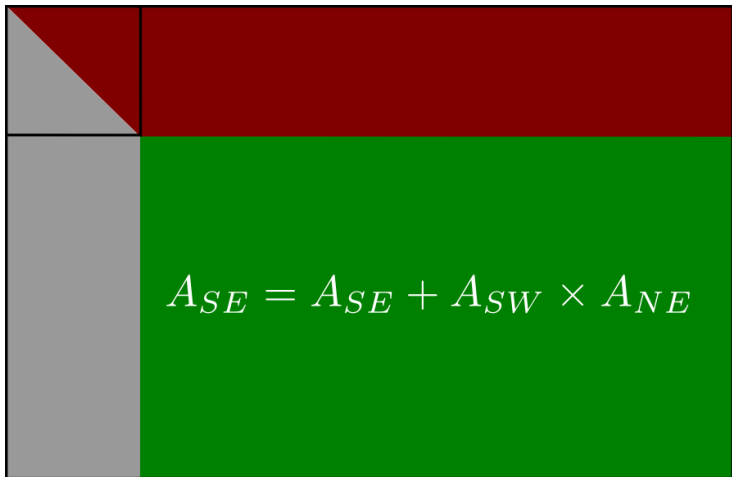
Visualisation





Visualisation

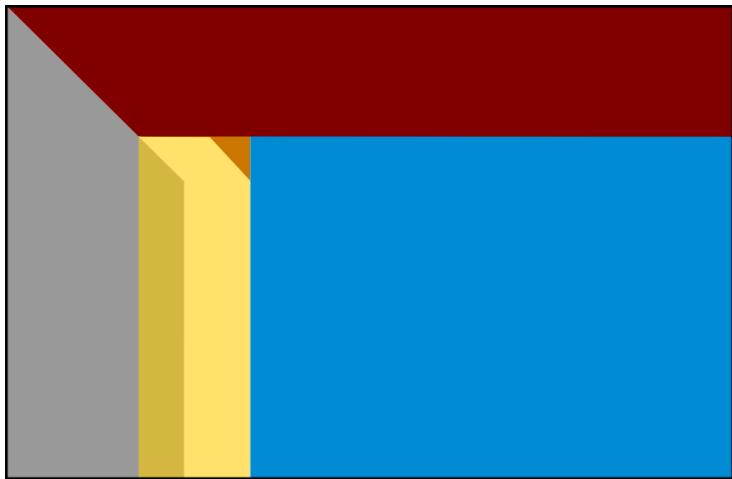




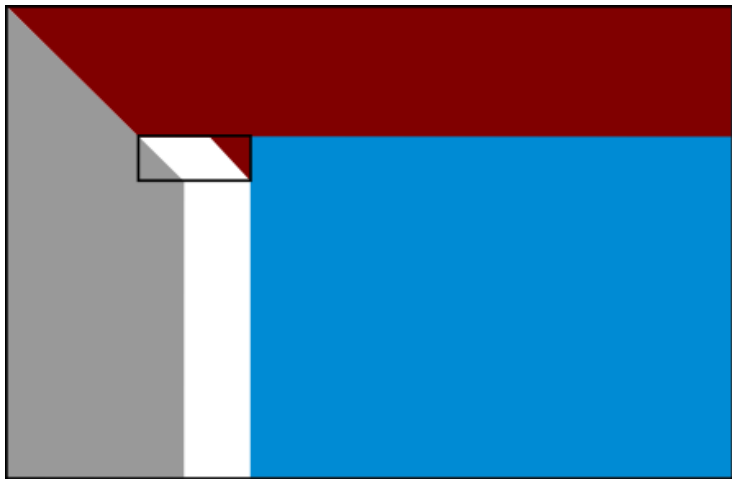
Visualisation



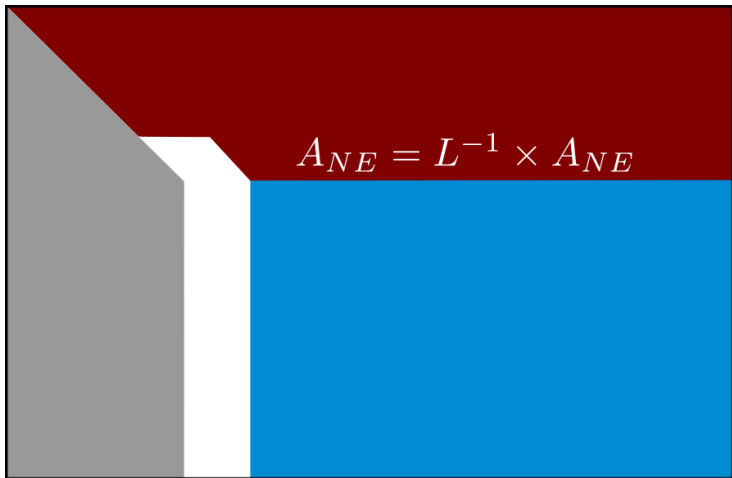
Visualisation



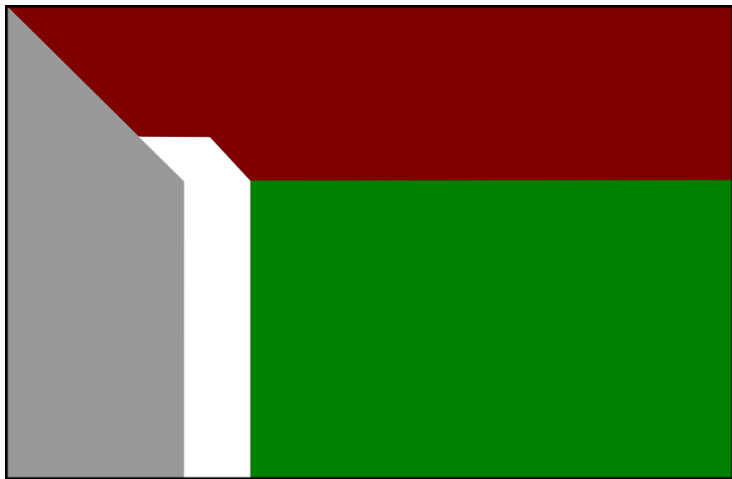
Visualisation



Visualisation



Visualisation



Visualisation



Results: Reduced Row Echelon Form

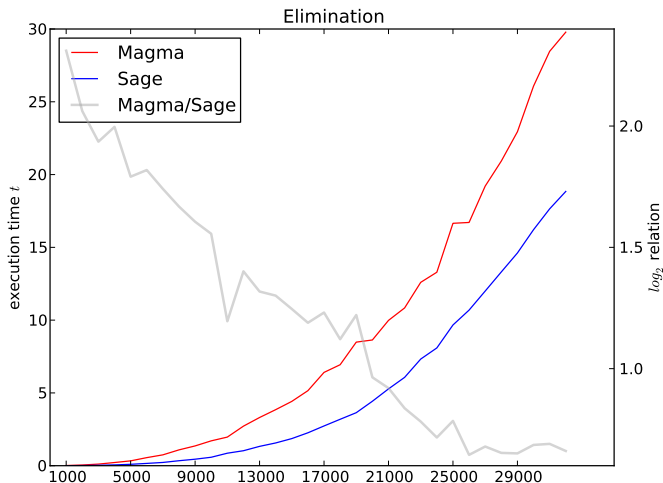


Figure: 2.66 Ghz Intel i7, 4GB RAM

Results: Row Echelon Form

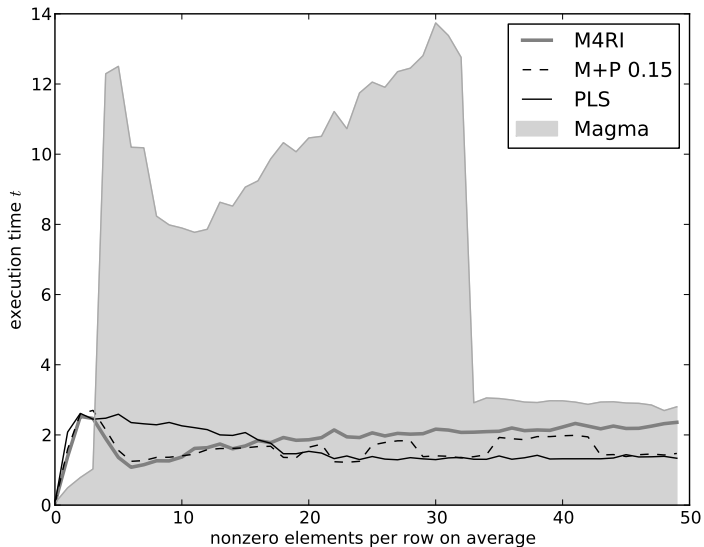
Using one core we can compute the echelon form of a $500,000 \times 500,000$ dense random matrix over \mathbb{F}_2 in

9711.42 seconds = 2.7 hours.

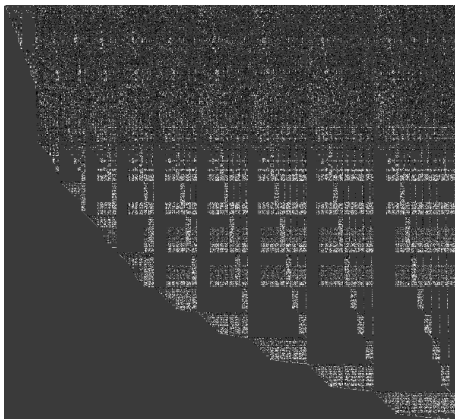
Using four cores decomposition we can compute the echelon form of a random dense $500,000 \times 500,000$ matrix in

3806.28 seconds = 1.05 hours.

Work-in-Progress: Sensitivity to Sparsity



Work-in-Progress: Gröbner Basis Linear Algebra



Problem	Matrix Dimension	Density	64-bit Debian/GNU Linux, 2.6Ghz Opteron)				
			Magma 2.15-10	M4RI 20100324	PLS 20100324	M+P 0.15 20100429	M+P 0.20 20100429
HFE 25	12,307 × 13,508	0.076	4.57s	3.28s	3.45s	3.03s	3.21s
HFE 30	19,907 × 29,323	0.067	33.21s	23.72s	25.42s	23.84s	25.09s
HFE 35	29,969 × 55,800	0.059	278.58s	126.08s	159.72s	154.62s	119.44s
MXL	26,075 × 26,407	0.185	76.81s	23.03s	19.04s	17.91s	18.00s

Work-in-Progress: Multi-core Support

